



27.09.2023

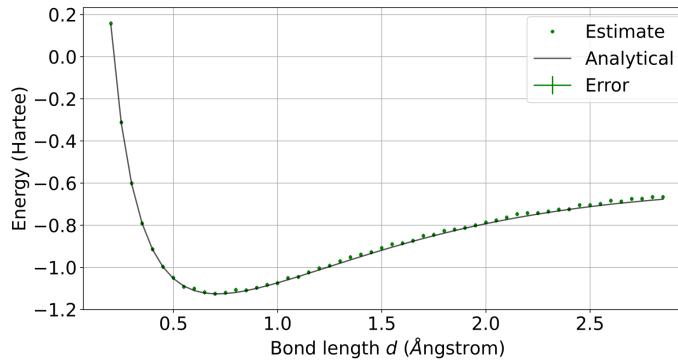
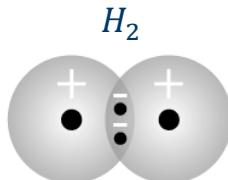
Adaptively measuring quantum expectation values using the empirical Bernstein stopping rule

Supervisor: Prof. Martin Kliesch
Ugur Tepe

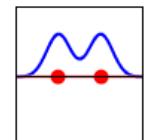
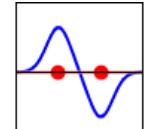
Why Quantum Computing

■ Physical Problems

1. Quantum Chemistry
2. Condensed Matter
3. ...



"Nature isn't classical, dammit,
and if you want to make a
simulation of nature, you'd better
make it quantum mechanical [...]"

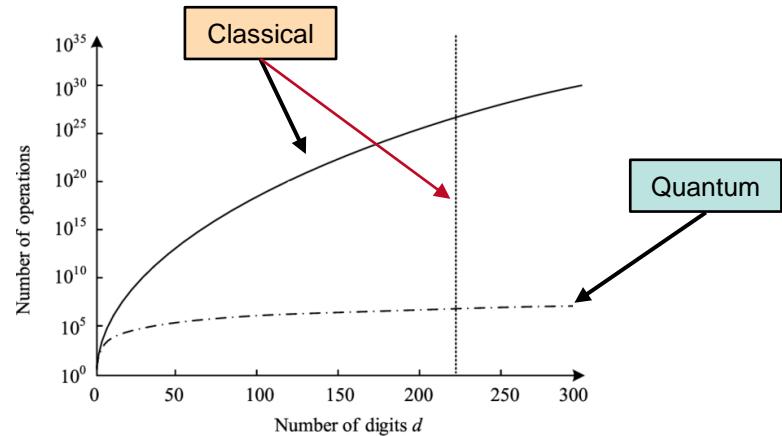


Why Quantum Computing

■ Mathematical Problems

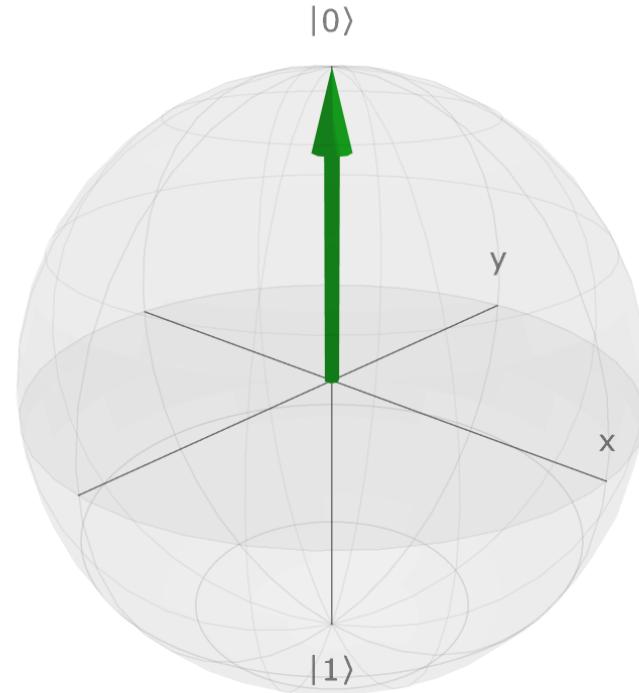
1. Solve System of Equations
2. Combinatorial Problems
3. Finding Prime Factors (Shor's algorithm)
4. ...

$$1517 = 41 * 37$$



Qubits

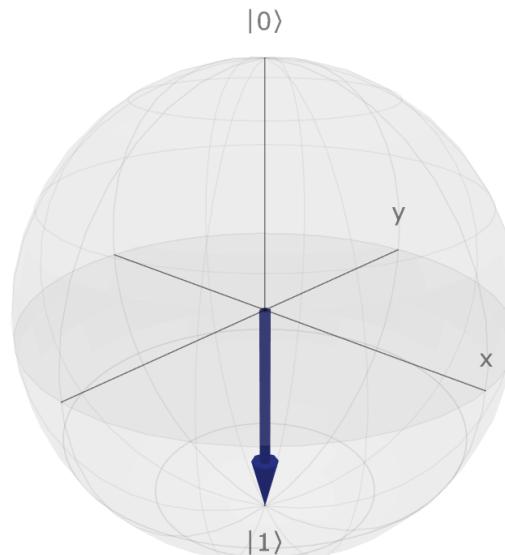
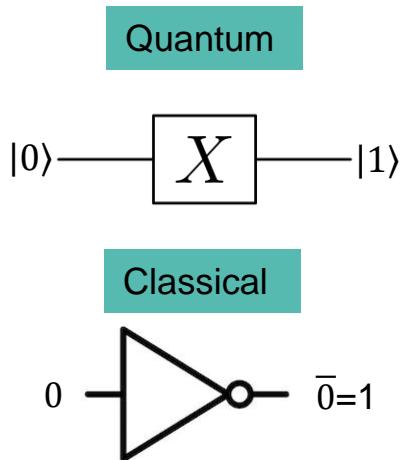
- Basis: $|0\rangle = \begin{bmatrix} 1 \\ 0 \end{bmatrix}$, $|1\rangle = \begin{bmatrix} 0 \\ 1 \end{bmatrix}$
- $|\psi\rangle = \alpha|0\rangle + \beta|1\rangle$
- Combine physical systems
→ Tensor product \otimes
- $|0\rangle \otimes |0\rangle = |00\rangle = \begin{bmatrix} 1 \\ 0 \end{bmatrix} \otimes \begin{bmatrix} 1 \\ 0 \end{bmatrix} = \begin{bmatrix} 1 \\ 0 \\ 0 \\ 0 \end{bmatrix}$
- $|01\rangle, |10\rangle, |11\rangle$ analogous



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Quantum Computing

Quantum Gates: X-gate



- $|0\rangle = \begin{bmatrix} 1 \\ 0 \end{bmatrix}$
- $|1\rangle = \begin{bmatrix} 0 \\ 1 \end{bmatrix}$
- X-gate:
 $\begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$

$$X|0\rangle = |1\rangle$$
$$X|1\rangle = |0\rangle$$

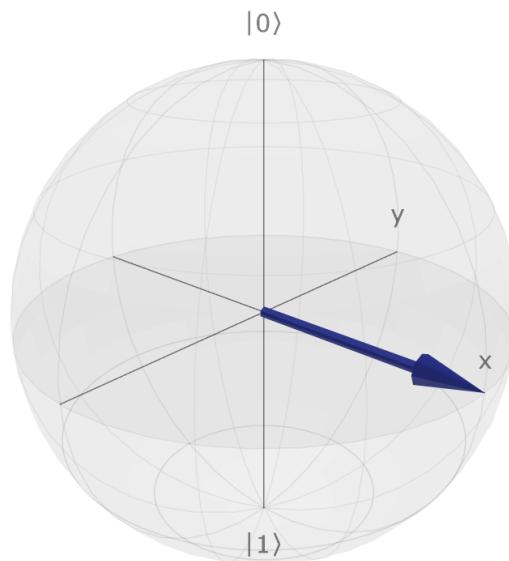
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Quantum Computing

Quantum Gates: R_y -gate

Quantum

$$|0\rangle \xrightarrow{R_Y(\theta)} \frac{|0\rangle + |1\rangle}{\sqrt{2}}$$



- $|0\rangle = \begin{bmatrix} 1 \\ 0 \end{bmatrix}$
- $|1\rangle = \begin{bmatrix} 0 \\ 1 \end{bmatrix}$
- $R_y(\theta) = \begin{bmatrix} \cos \theta/2 & -\sin \theta/2 \\ \sin \theta & \cos \theta/2 \end{bmatrix}$

$$R_y\left(\frac{\pi}{2}\right)|0\rangle = \frac{|0\rangle + |1\rangle}{\sqrt{2}}$$
$$R_y\left(\frac{\pi}{2}\right)|1\rangle = \frac{|0\rangle - |1\rangle}{\sqrt{2}}$$

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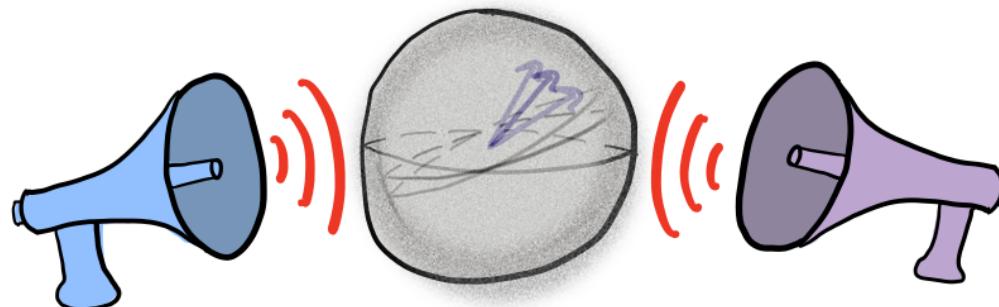
Limitation of current quantum computers/ hardware

Noise



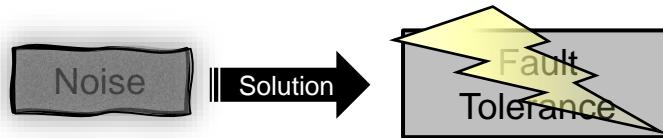
Quantum Computing

Limitation of current quantum computers/ hardware

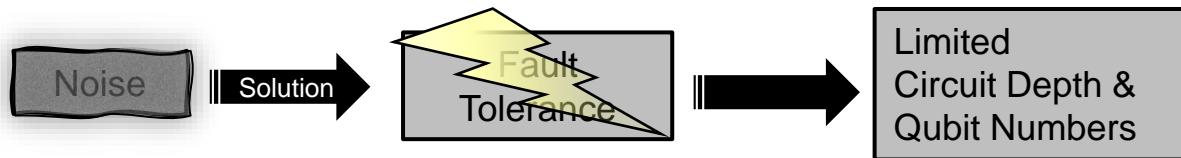


Quantum Computing

Limitation of current quantum computers/ hardware

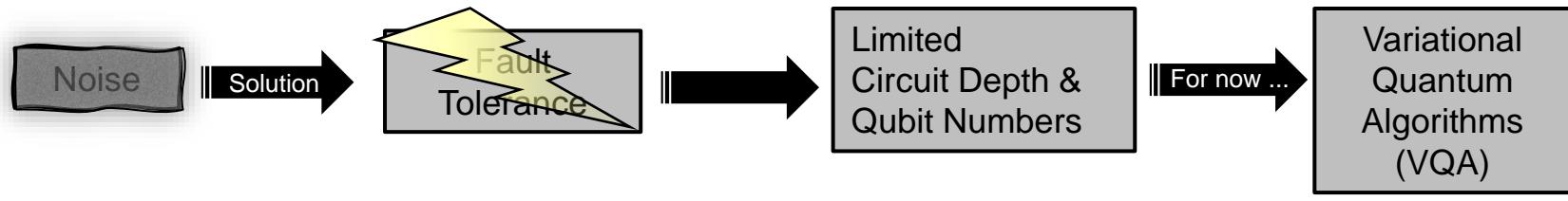


Limitation of current quantum computers/ hardware

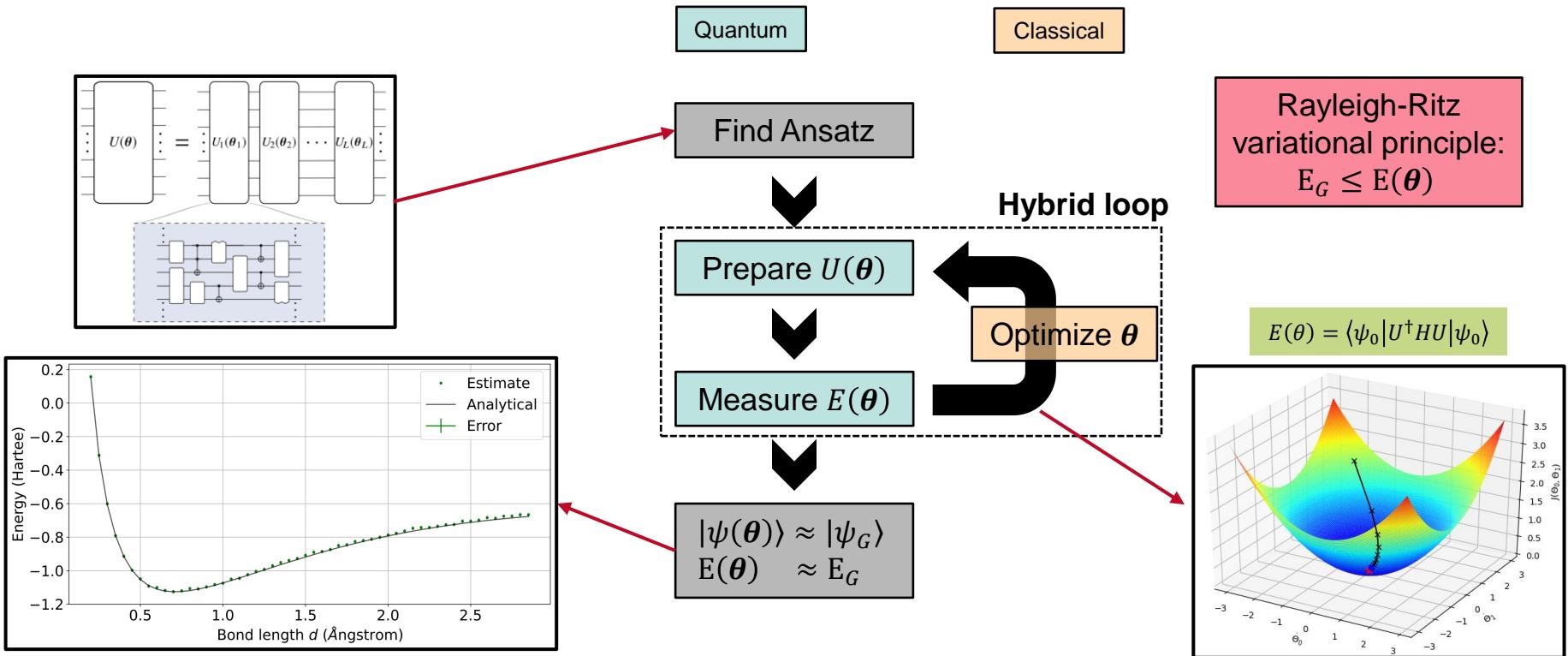


Quantum Computing

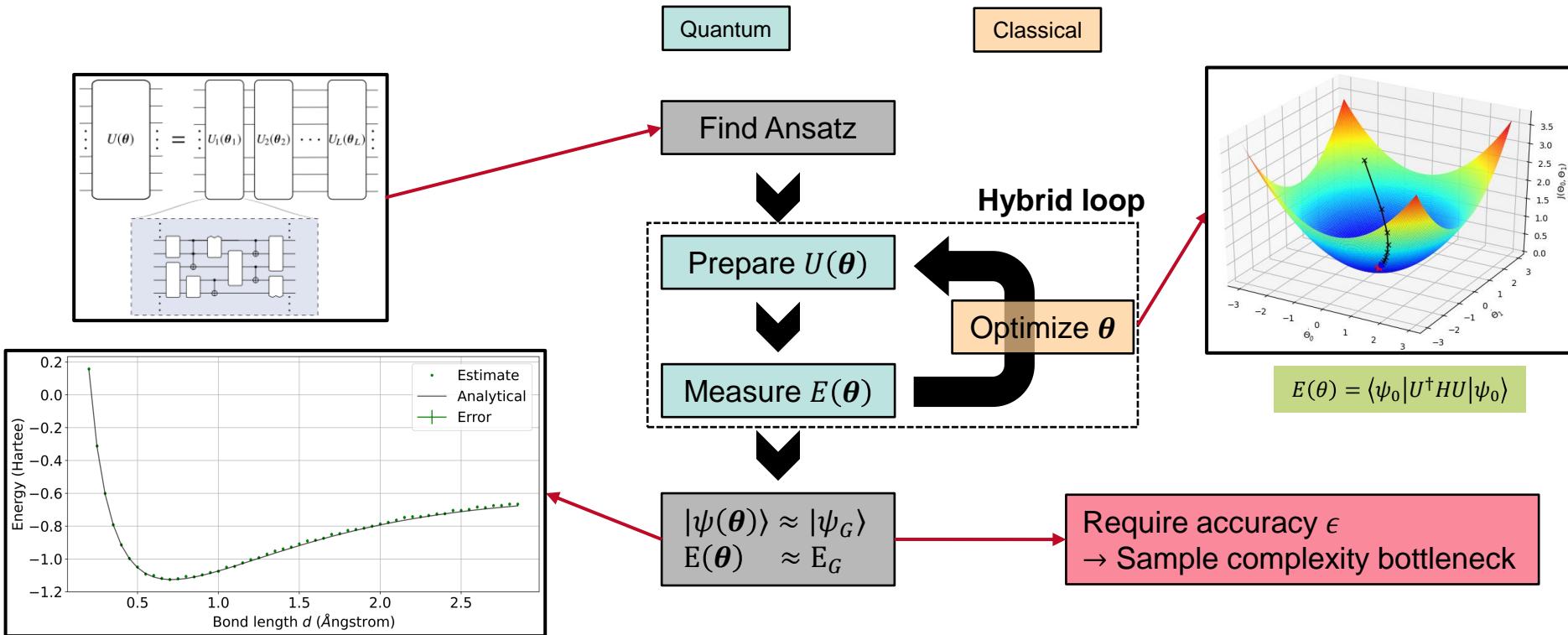
Limitation of current quantum computers/ hardware



Schematic Overview of VQA



Schematic Overview of VQA



Empirical Bernstein Stopping (EBS)

An adaptive sampling algorithm

Bernstein's Inequality

Independent and identically distributed random variable

$$\mathbb{P}[|\bar{X}_t - \mu| \geq \epsilon] \leq \exp\left[-\frac{\frac{1}{2}(t\epsilon)^2}{\Sigma^2 + \frac{1}{3}Rt\epsilon}\right] := \delta$$

- So called bound form

$$|\bar{X}_t - \mu| \leq \sqrt{\frac{2 \Sigma^2 \ln(2/\delta)}{t}} + \frac{R \ln(2/\delta)}{3t}$$

- $X_1 \dots X_t \triangleq$ i.i.d random variables
- $a \leq X_i \leq b$
- $\epsilon \in R^+$
- $R \triangleq$ Range

- $\bar{X}_t = \frac{1}{t} \sum_{i=1}^t X_i$
- $\Sigma^2 = \sum_{i=1}^t \sigma_i^2$

Bernstein's Inequality

$$\mathbb{P}[|\bar{X}_t - \mu| \geq \epsilon] \leq \exp\left[-\frac{\frac{1}{2}(t\epsilon)^2}{\Sigma^2 + \frac{1}{3}Rt\epsilon}\right] := \delta$$

- So called bound form

$$|\bar{X}_t - \mu| \leq \sqrt{\frac{2 \Sigma^2 \ln(2/\delta)}{t}} + \frac{R \ln(2/\delta)}{3t}$$

- Variance usually unknown !

Independent and identically distributed random variable

- $X_1 \dots X_t \triangleq$ i.i.d random variables
- $a \leq X_i \leq b$
- $\epsilon \in R^+$
- $R \triangleq$ Range

- $\bar{X}_t = \frac{1}{t} \sum_{i=1}^t X_i$
- $\Sigma^2 = \sum_{i=1}^t \sigma_i^2$

Empirical Bernstein Bound

Variance Σ^2 Replace Empirical variance \bar{V}_t

■ Empirical Bernstein Bound¹

$$|\bar{X}_t - \mu| \leq \sqrt{\frac{2\bar{V}_t \ln(3/\delta)}{t}} + \frac{3R \ln(3/\delta)}{t}$$

$\underbrace{\phantom{\sqrt{\frac{2\bar{V}_t \ln(3/\delta)}{t}}}}_{\Theta\left(\frac{1}{\sqrt{t}}\right)}$ $\underbrace{\phantom{\frac{3R \ln(3/\delta)}{t}}}_{\Theta\left(\frac{1}{t}\right)}$

■ Ideally: $\bar{V}_t \ll R^2$!

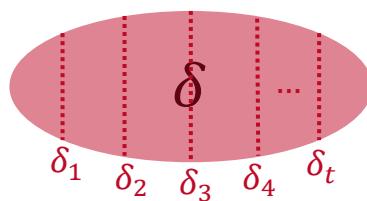
Independent and identically distributed random variable

- $X_1 \dots X_t \triangleq$ i.i.d random variables
- $a \leq X_i \leq b$
- $\epsilon \in R^+$
- $R \triangleq$ Range

- $\bar{X}_t = \frac{1}{t} \sum_{i=1}^t X_i$
- $\bar{V}_t = \frac{1}{t} \sum_{i=1}^t (X_i - \bar{X}_t)^2$

Empirical Bernstein Stopping (EBS)²

Pseudo Code Implementation



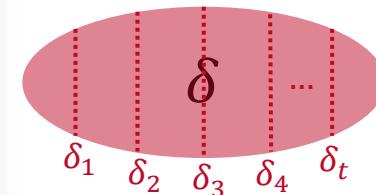
- $X_1 \dots X_t \triangleq$ i.i.d random variables
- $\epsilon \in R^+$
- $1 - \delta \triangleq$ Confidence
- $R \triangleq$ Range

$$\begin{aligned} & \bullet \quad \bar{X}_t = \frac{1}{t} \sum_{i=1}^t X_i \\ & \bullet \quad c(t) = \sqrt{\frac{2\bar{V}_t \ln(3/\delta_t)}{t}} + \frac{3R \ln(3/\delta_t)}{t} \end{aligned}$$

Empirical Bernstein Stopping (EBS)²

Pseudo Code Implementation

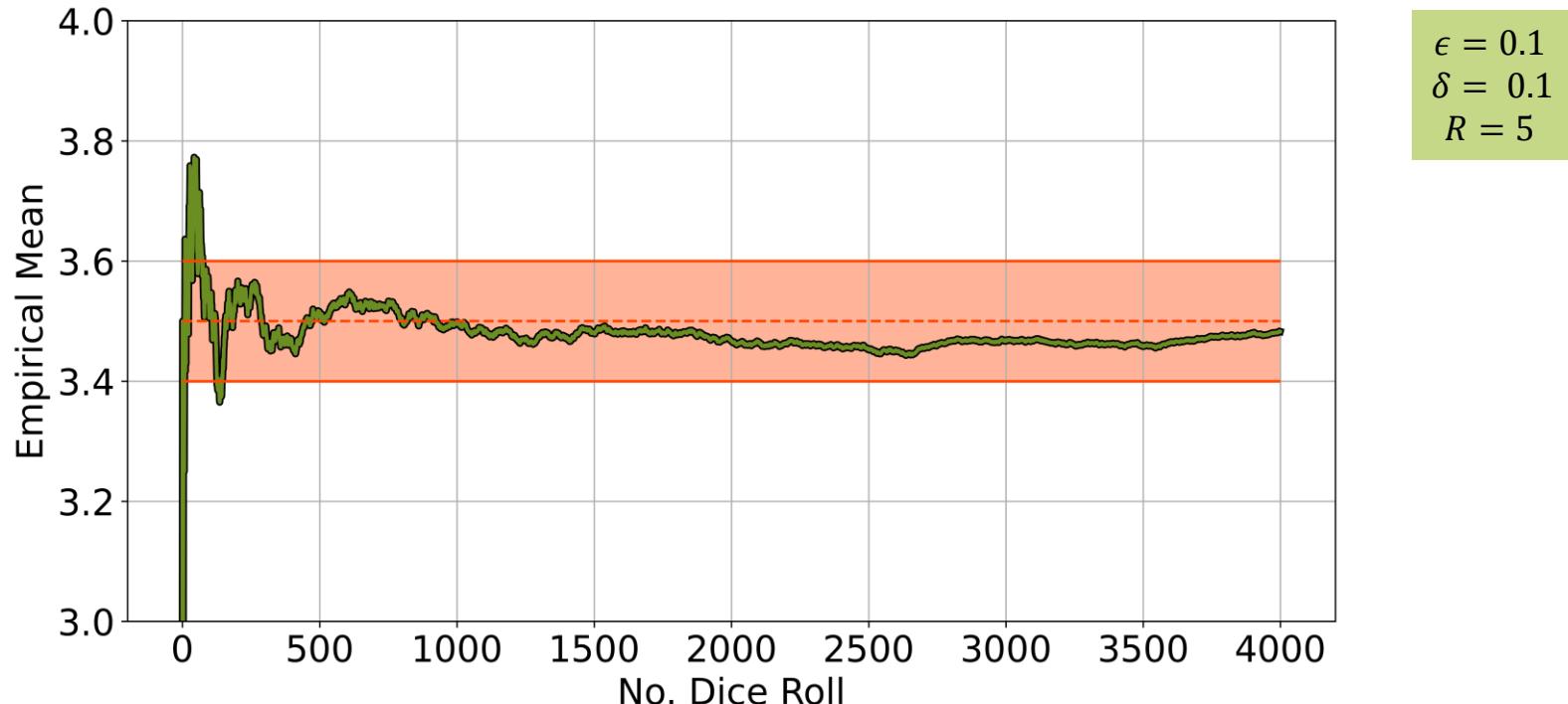
```
while ct > eps:  
    sample(X)  
    update(mean, variance)  
    update(ct)  
  
return mean
```



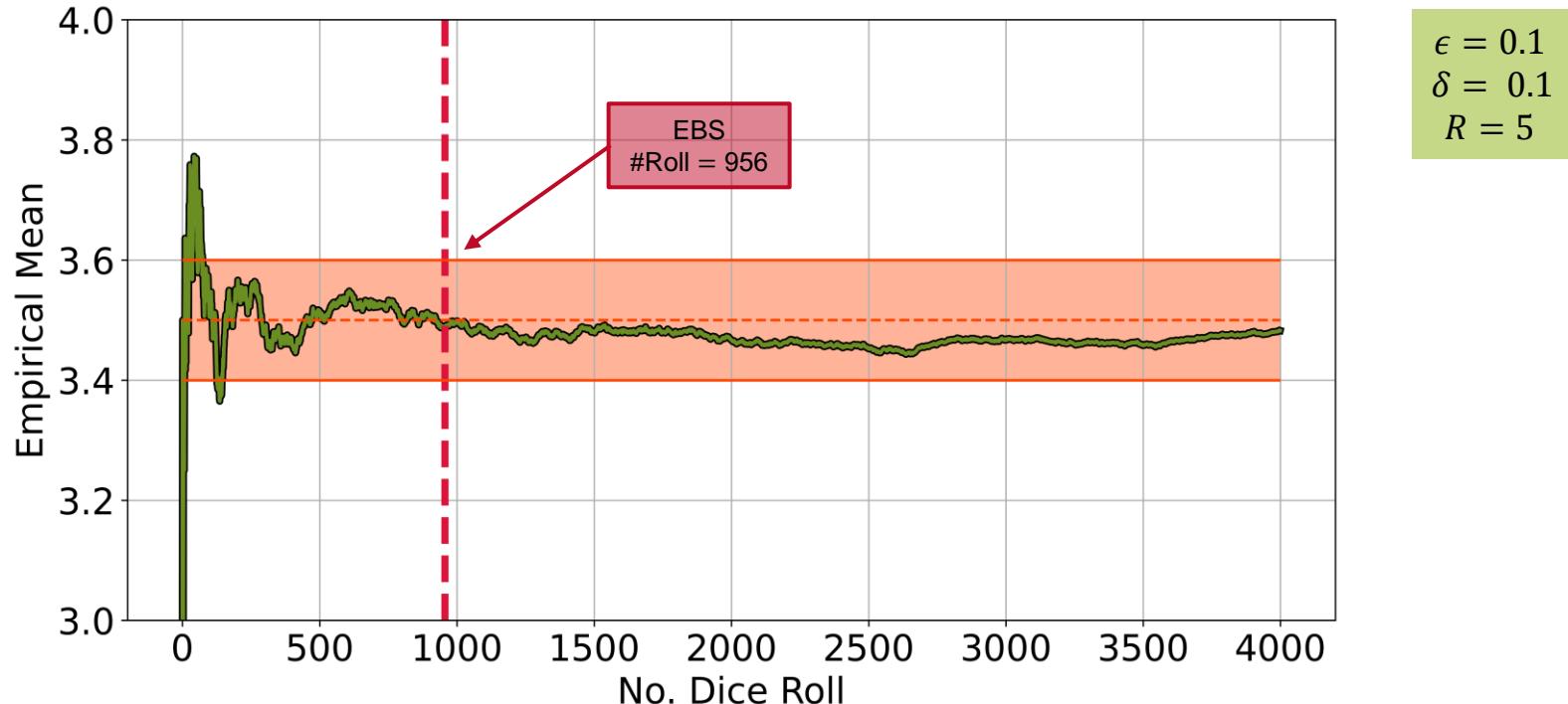
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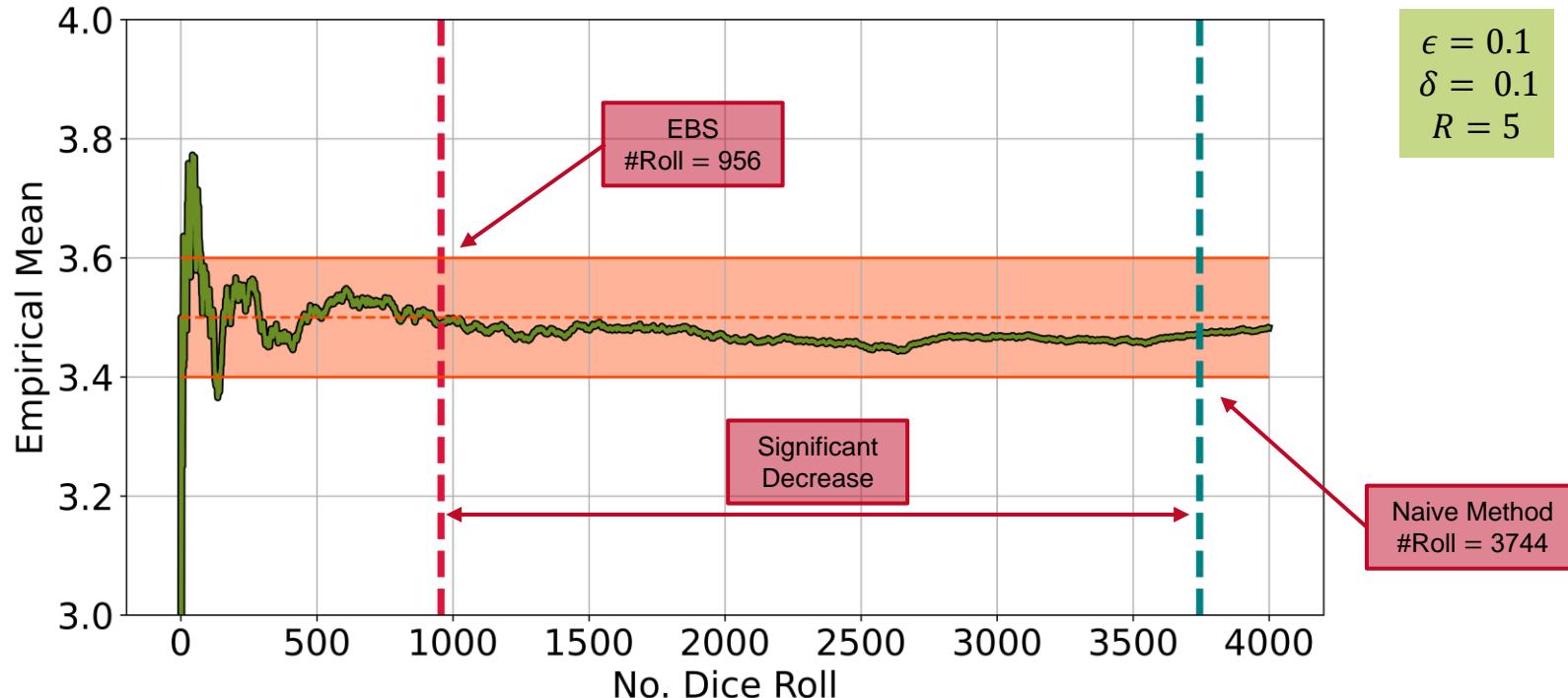
Empirical Mean of a Dice: EBS Algorithm



Empirical Mean of a Dice: EBS Algorithm



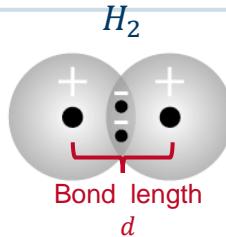
Empirical Mean of a Dice: EBS Algorithm



VQAs for solving the Electronic Structure Problem

Electronic Structure Problem

- Describing whole system



Kinetic Energy

$$H = - \sum_i \frac{\nabla_{R_i}^2}{2M_i} - \sum_i \frac{\nabla_{r_i}^2}{2m_i}$$

Nuclei
Electron

Potential Energy

$$\sum_{i,j} \frac{Z_i e^2}{|R_i - r_j|} + \sum_{i,j>i} \frac{Z_i Z_j e^2}{|R_i - R_j|} - \sum_{i,j>i} \frac{e^2}{|r_i - r_j|}$$

Electron - Nuclei
Nuclei - Nuclei
Electron - Electron

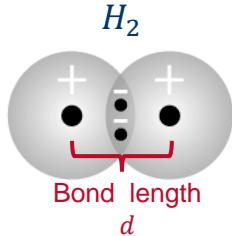
- $R_i \triangleq$ Nuclei position
- $r_i \triangleq$ Electron position
- $M_i \triangleq$ Nuclei mass
- $m_i \triangleq$ Electron mass
- $Z_i \triangleq$ Nuclei charge
- 1 Hartee \triangleq 27.2 eV

- Born-Oppenheimer approximation + second quantization

$$H = \sum_{p,q} h_{pq} a_p^\dagger a_q + \frac{1}{2} \sum_{p,q,r,s} h_{pqrs} a_p^\dagger a_q^\dagger a_r a_s$$

- $a^\dagger \triangleq$ creation operator
- $a \triangleq$ annihilation operator
- p, q, r, s : orbital label

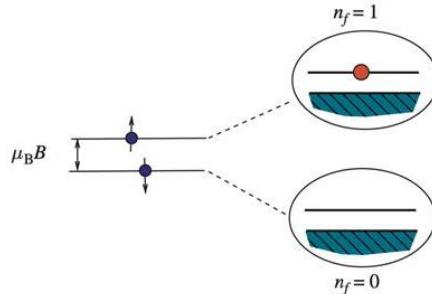
Electronic Structure Problem



$$H = \sum_{p,q} h_{pq} a_p^\dagger a_q + \frac{1}{2} \sum_{p,q,r,s} h_{pqrs} a_p^\dagger a_q^\dagger a_r a_s$$

- $R_i \triangleq$ Nuclei position
- $r_i \triangleq$ Electron position
- $M_i \triangleq$ Nuclei mass
- $m_i \triangleq$ Electron mass
- $Z_i \triangleq$ Nuclei charge
- 1 Hartee $\triangleq 27.2 \text{ eV}$

- How to represent this on a quantum computer ?
 - Jordan – Wigner mapping
 - ...
- Fermionic operators \rightarrow Pauli operators $\sum_i c_i P_i$



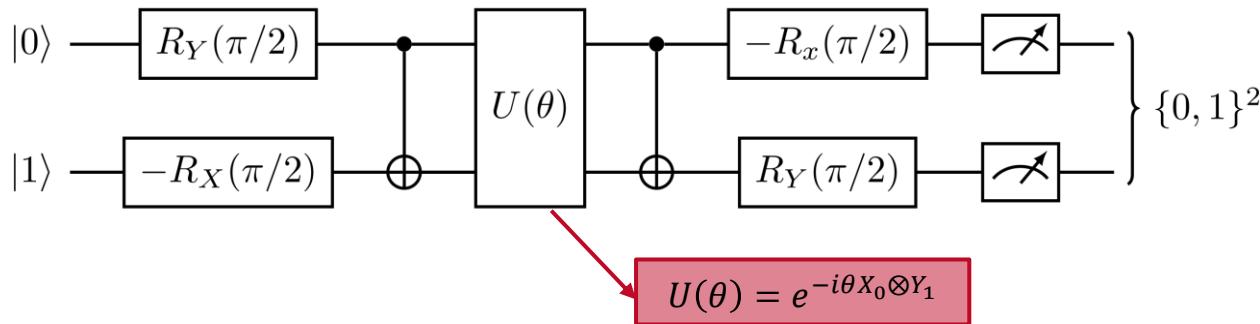
- $a^\dagger \triangleq$ creation operator
- $a \triangleq$ annihilation operator
- p, q, r, s : orbital label

Results: Total Energy of H_2

Total Energy of H_2

Overview

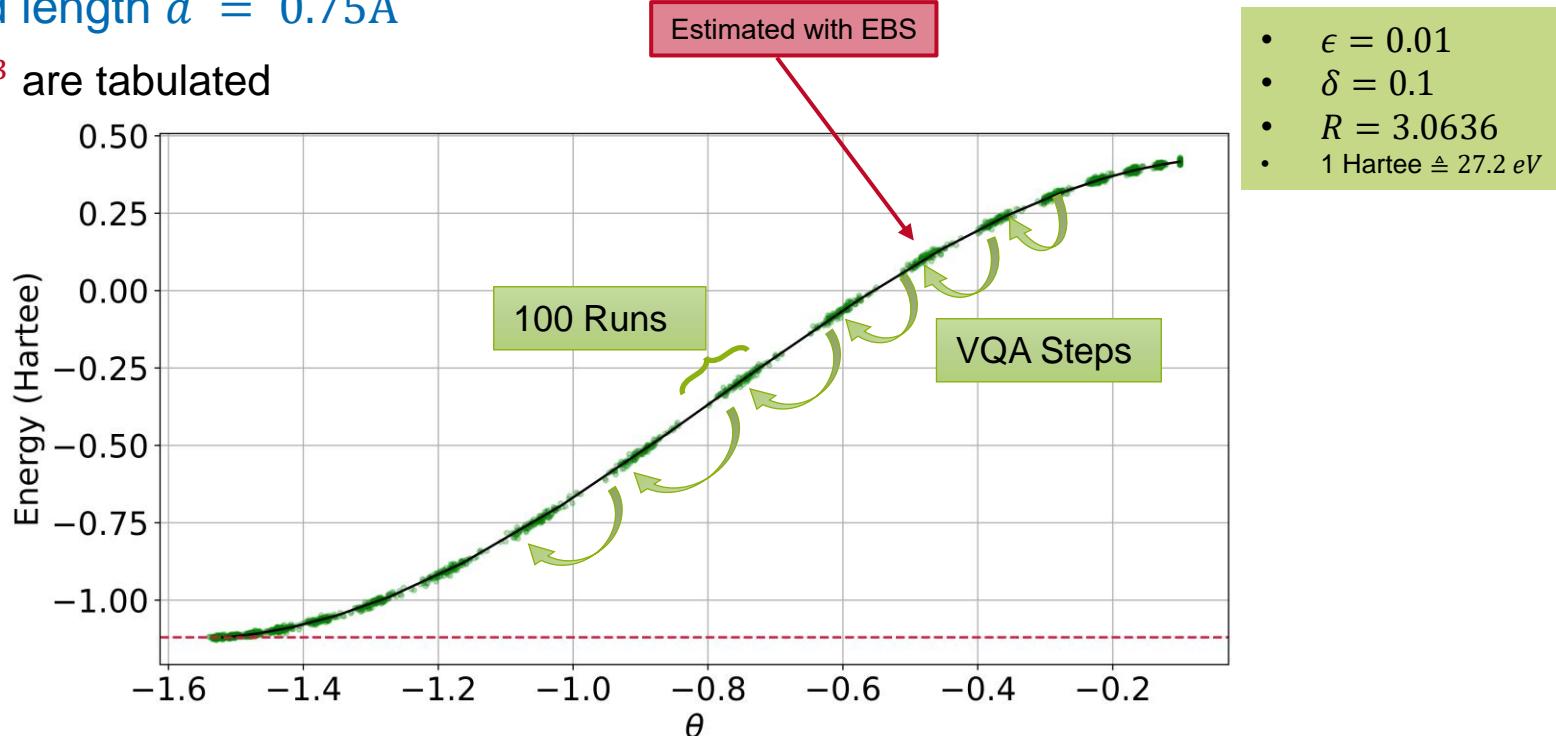
- $\hat{H}_{H_2}(d) = g_1(d)\mathbb{I} + g_2(d)Z_0 + g_3(d)Z_1 + g_4(d)Z_0 \otimes Z_1 + g_5(d)Y_0 \otimes Y_1 + g_6(d)X_0 \otimes X_1$ ³
- Guaranteed accuracy ϵ with EBS
- Ansatz for parametrised Circuit: ³



Ground State Energy $H_2(d)$: VQA

Fixed bond length $d = 0.75\text{\AA}$

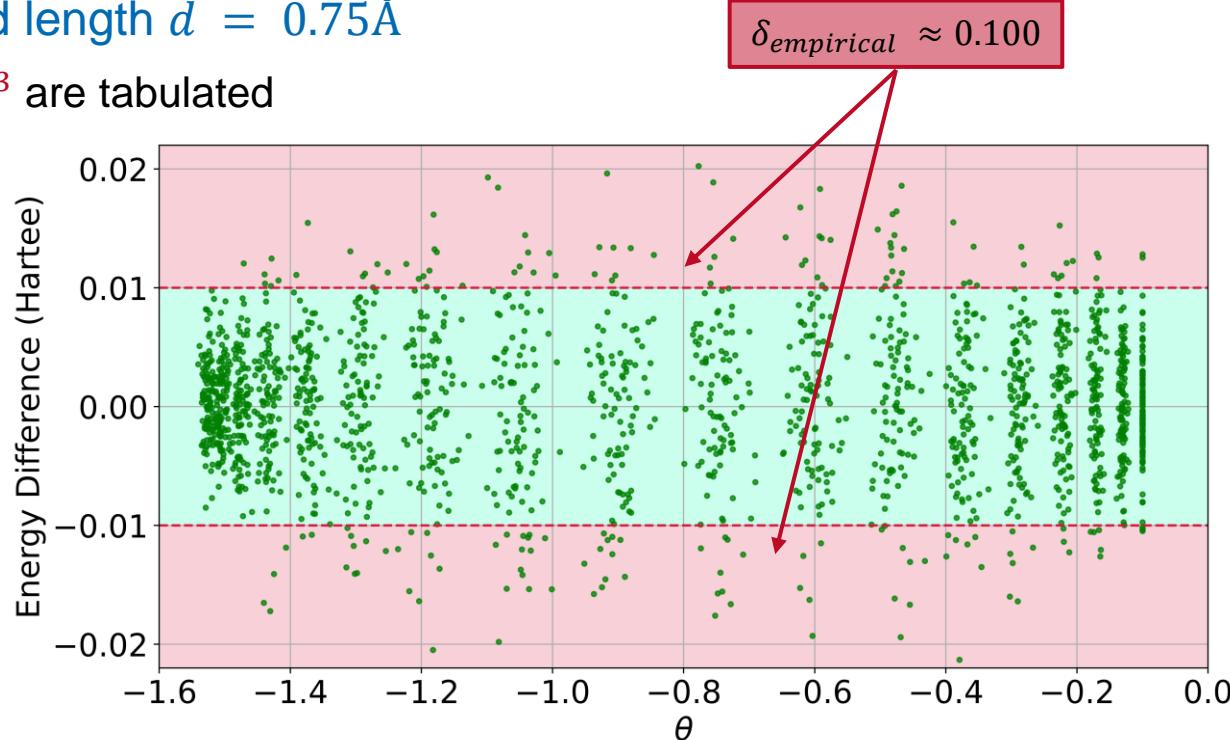
- $d \rightarrow \{g_i\}^3$ are tabulated



Ground State Energy $H_2(d)$: VQA

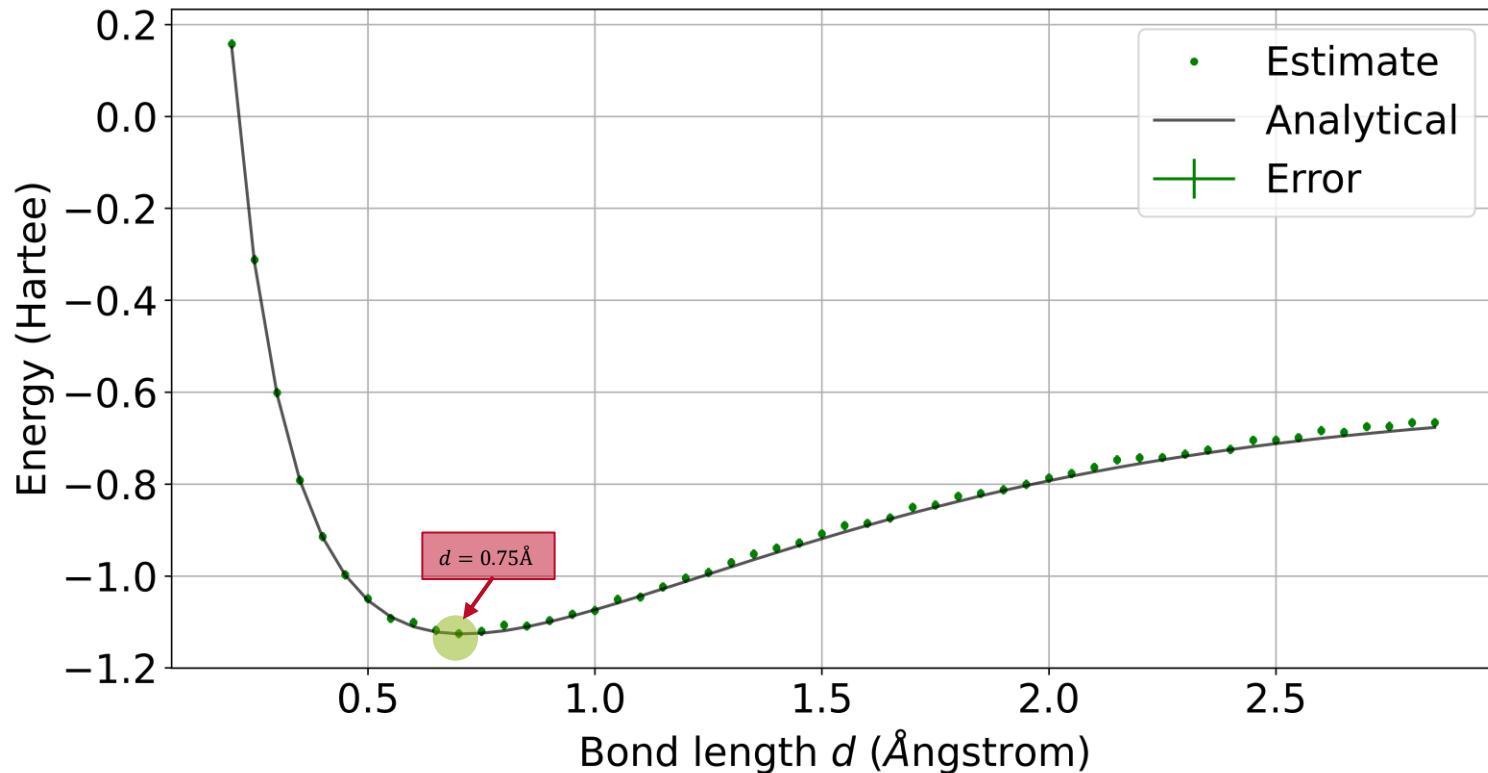
Fixed bond length $d = 0.75\text{\AA}$

- $d \rightarrow \{g_i\}^3$ are tabulated



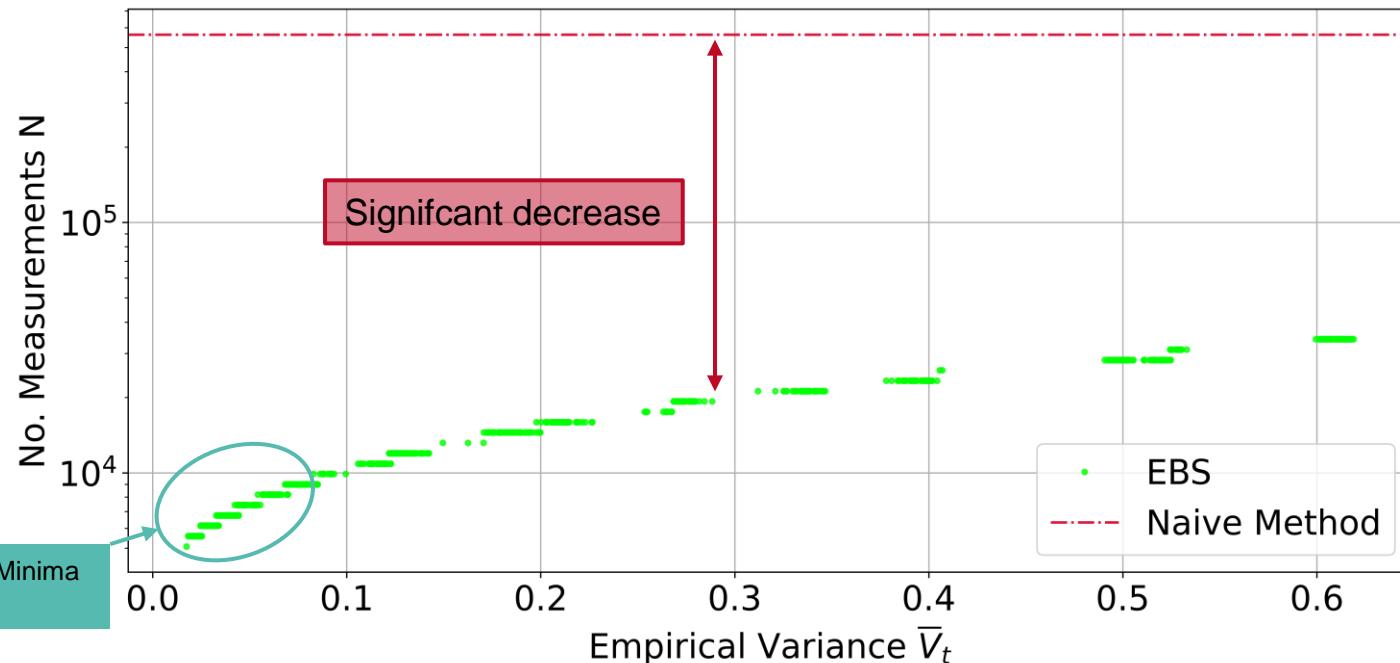
- $\epsilon = 0.01$
- $\delta = 0.1$
- $R = 3.0636$
- 1 Hartree $\triangleq 27.2 \text{ eV}$

Total Energy of H_2



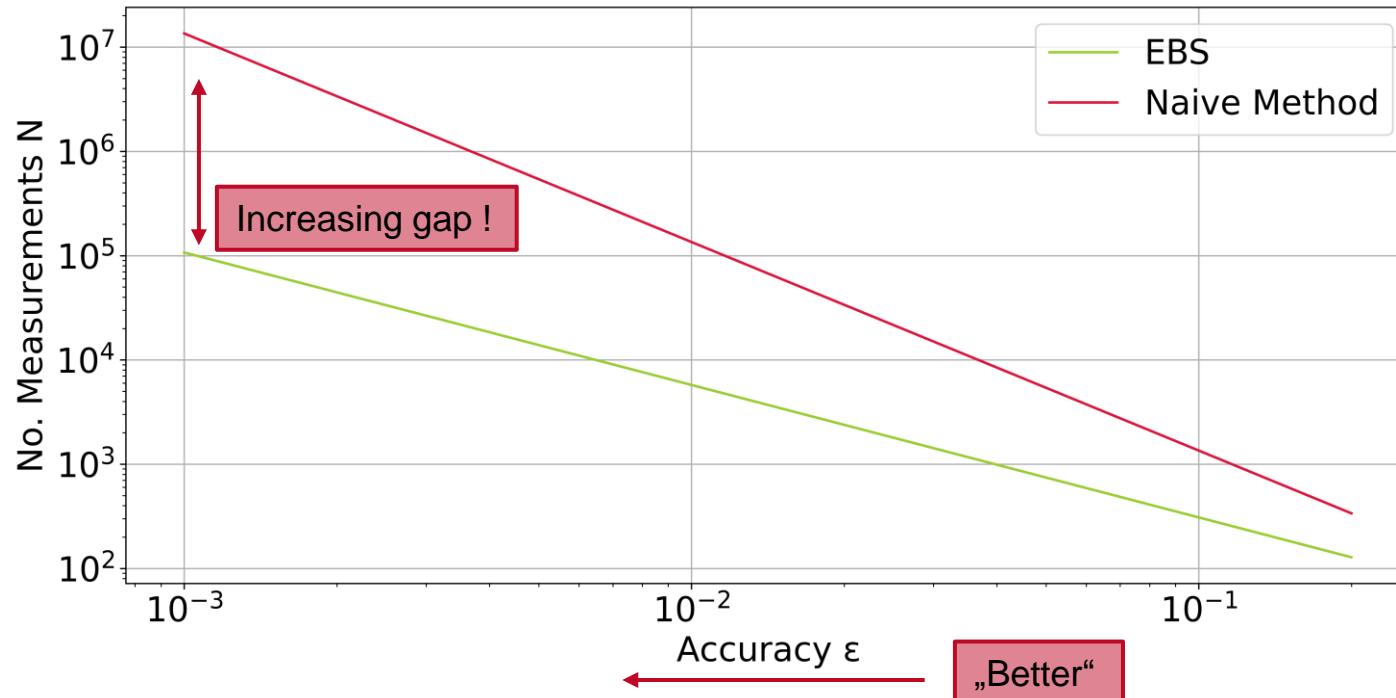
Reducing the measurement effort

Empirical Variance \bar{V}_t vs Number of Measurements N



Reducing the measurement effort

Accuracy ϵ vs Number of Measurements N



Summary & Outlook

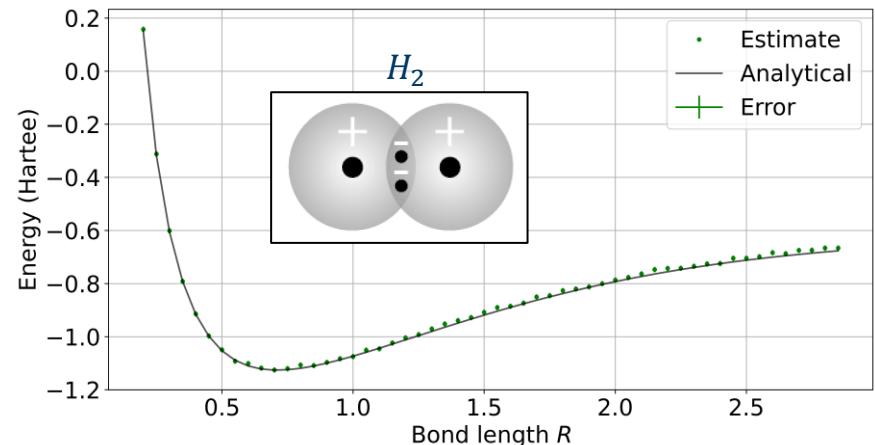
Summary & Outlook

Summary:

- **Significant reduction of VQA measurement effort**
- **Guaranteed** confidence $1 - \delta$ and accuracy ϵ
However, ...
- **Finding** suitable **ansatz** is **hard** →
EBS advantage may not be a pronounced
- H_2 is **minimal** example
- Optimal setup → **no noise** simulation
- **Energy** estimator depends on
measurement strategy

Outlook

- Different applications: e.g., state verifications, etc...
- Test on **real** quantum hardware
- **Improvements** to EBS: Adapt **beyond real-valued random variables**



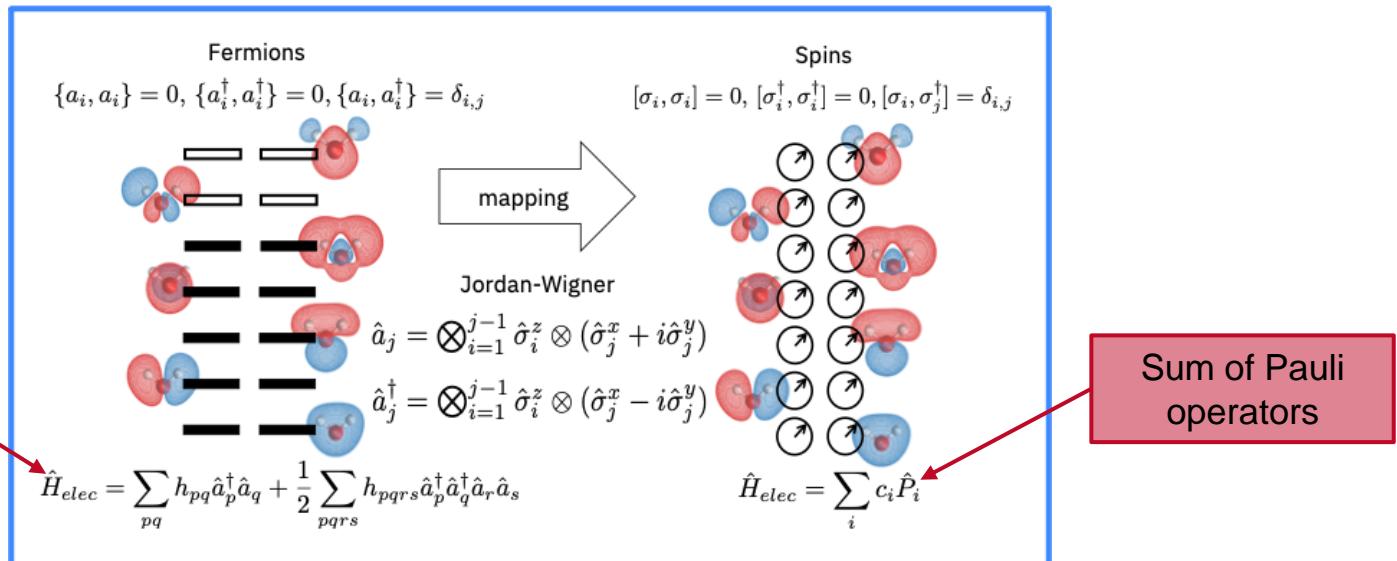
1. JY. Audibert et. al., *Tuning bandit algorithms in stochastic environments* (2007)
2. V. Mnih, *Efficient stopping rules* (2008)
3. P. J. O'Malley, *Scalable quantum simulation of molecular energies* (2016)
4. J. Preskill, *Quantum computing in the NISQ era and beyond* (2018)
5. A. Peruzzo, *A variational eigenvalue solver on a photonic quantum processor* (2018)
6. J. K. Blitzstein and J. Hwang, *Introduction to probability* (2015)
7. M. Kliesch, *Characterization, certification, and validation of quantum systems* (2020)
8. M. A. Nielsen, *Quantum computation and quantum information* (2010)
9. Smite-Meister, *Bloch sphere, a geometrical representation of a two-level quantum system* (2009)
10. L. Zhu, *Optimizing shot assignment in variational quantum eigensolver measurement* (2023)
11. Molecular orbit H2 <https://commons.wikimedia.org/wiki/File:H2OrbitalsAnimation.gif> 25.09.2023 (20:15)
12. Feynman picture <https://www.britannica.com/biography/Richard-Feynman> 25.09.2023 (20:15)
13. H2 molecule <https://byjus.com/chemistry/hydrogen-gas/> 25.09.2023 (20:15)

Appendix: Jordan – Wigner Mapping

Electronic Structure Problem

Jordan-Wigner Mapping

- 1 electron orbital + spin → 2 qubits



Measurement

Energy curve of H_2

How to measure?

- $\hat{H}_{H_2}(R) = g_1(R)\mathbb{I} + g_2(R)Z_0 + g_3(R)Z_1 + g_4(R)Z_0 \otimes Z_1 + g_5(R)Y_0 \otimes Y_1 + g_6(R)X_0 \otimes X_1$

\hat{H}_i	Measurement Basis	$E(R)$ $+g_0E_0$
Z_0		g_2E_2
Z_1	$\left. \begin{array}{c} \\ \end{array} \right\}$	g_3E_3
$Z_0 \otimes Z_1$	$\left. \begin{array}{c} \\ \end{array} \right\}$	g_4E_4
$Y_0 \otimes Y_1$	$\left. \begin{array}{c} \\ \end{array} \right\}$	g_5E_5
$X_0 \otimes X_1$	$\left. \begin{array}{c} \\ \end{array} \right\}$	g_6E_6

Energy curve of H_2

How to measure?

■ $\hat{H}_{H_2}(R) = g_1(R)\mathbb{I} + g_2(R)Z_0 + g_3(R)Z_1 + g_4(R)Z_0 \otimes Z_1 + g_5(R)Y_0 \otimes Y_1 + g_6(R)X_0 \otimes X_1$

\hat{H}_i	Measurement Basis	$E(R) + g_0 E_0$
Z_0		$g_2 E_2$
Z_1	 z	$g_3 E_3$
$Z_0 \otimes Z_1$		$g_4 E_4$
$Y_0 \otimes Y_1$	 Y	$g_5 E_5$
$X_0 \otimes X_1$	 X	$g_6 E_6$

- 3 x measurement
→ estimate
- General optimal strategy **unclear**

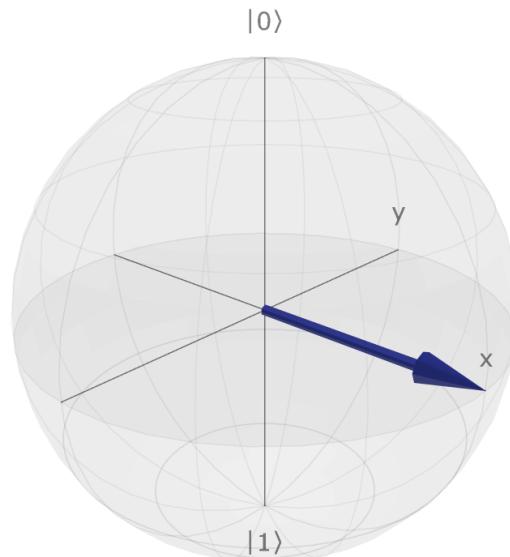
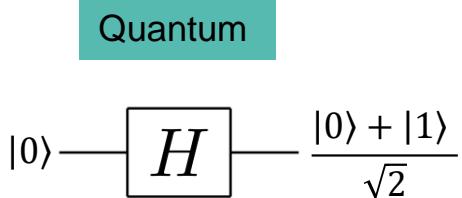
Quantum Computing

Physical Qubit Implementations

Physical Support	Information	$ 0\rangle$	$ 1\rangle$
Photons	Polarization	\leftrightarrow	\updownarrow
	Number	Vacuum	Single Photon
Electron	Spin	\uparrow	\downarrow
	Number	No Electron	One Electron
...

Quantum Computing

Quantum Gates

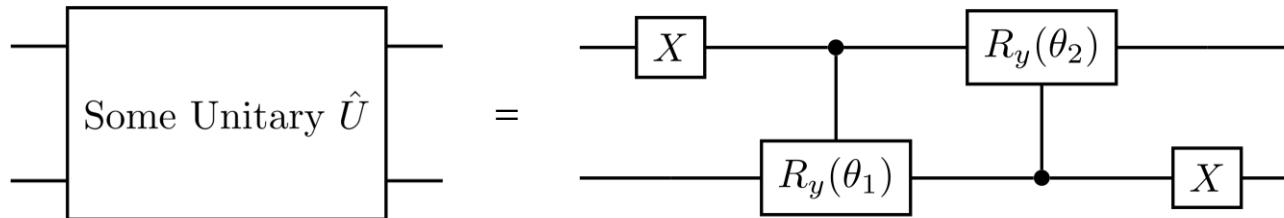


- $|0\rangle = \begin{bmatrix} 1 \\ 0 \end{bmatrix}$
- $|1\rangle = \begin{bmatrix} 0 \\ 1 \end{bmatrix}$
- Hadamard -gate= $\frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix}$

$$H|0\rangle = \frac{1}{\sqrt{2}}|0\rangle + \frac{1}{\sqrt{2}}|1\rangle$$
$$H|1\rangle = \frac{1}{\sqrt{2}}|0\rangle - \frac{1}{\sqrt{2}}|1\rangle$$

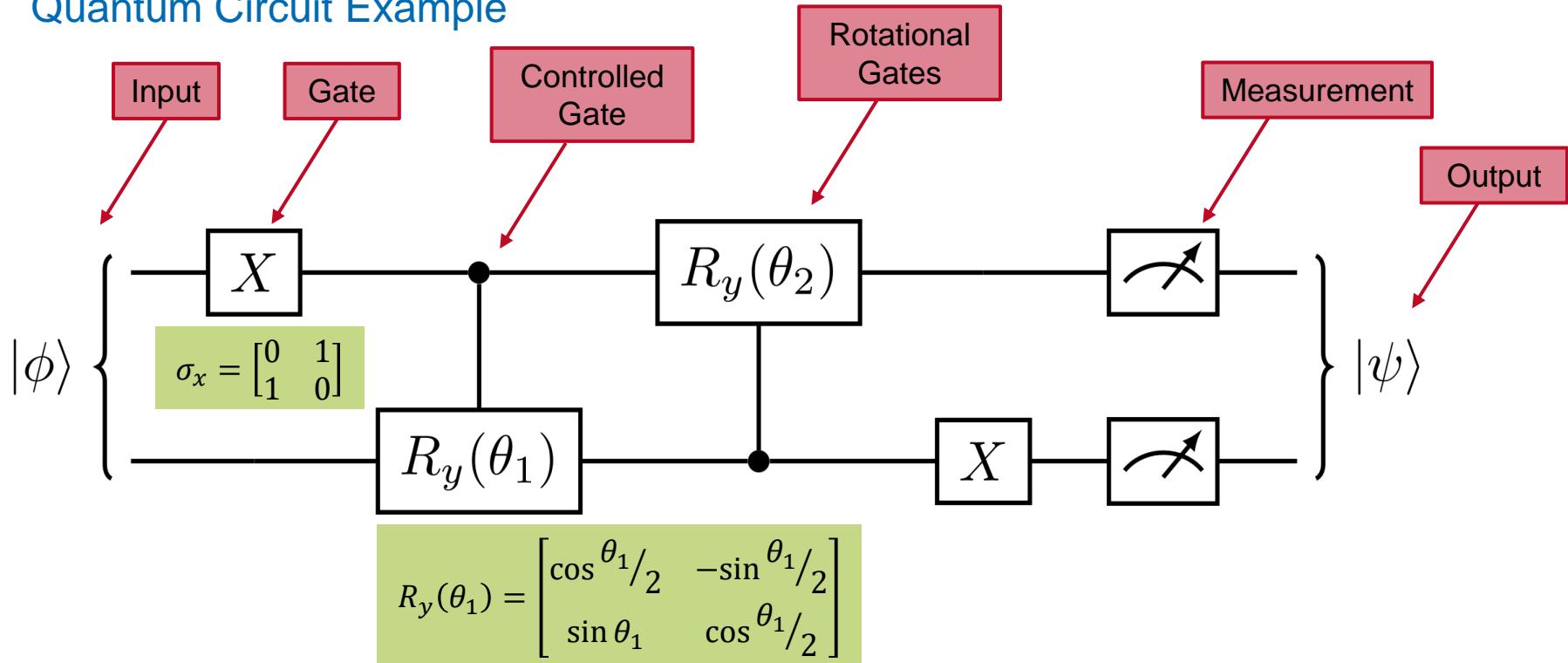
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- Certain sets of gates are universal
 - Every unitary operation can be built using these
- Any gate can be decomposed using these
 - {CNOT, all single qubit gates}¹
 - {CNOT, H , T }²
 - ...



Quantum Computing

Quantum Circuit Example



Tail Bounds

- A stopping algorithm terminates when condition is met, i.e.:

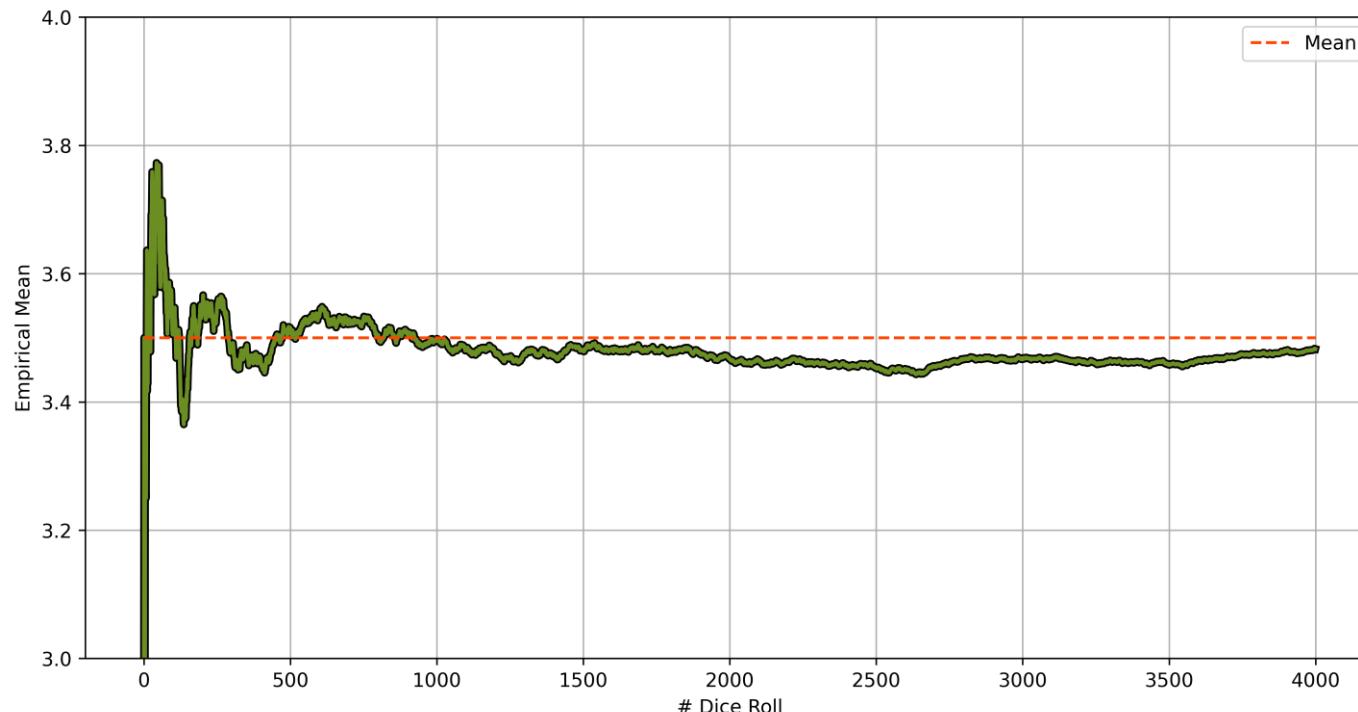
$$\mathbb{P}[|\hat{\mu} - \mu| \leq \epsilon] \geq 1 - \delta$$

- Estimate $\hat{\mu}$ is called an absolute (ϵ, δ) -estimate

- $\hat{\mu} \triangleq$ Estimate
- $\mu \triangleq$ Expected value
- $\epsilon \triangleq$ Error margin
- $1 - \delta \triangleq$ Confidence

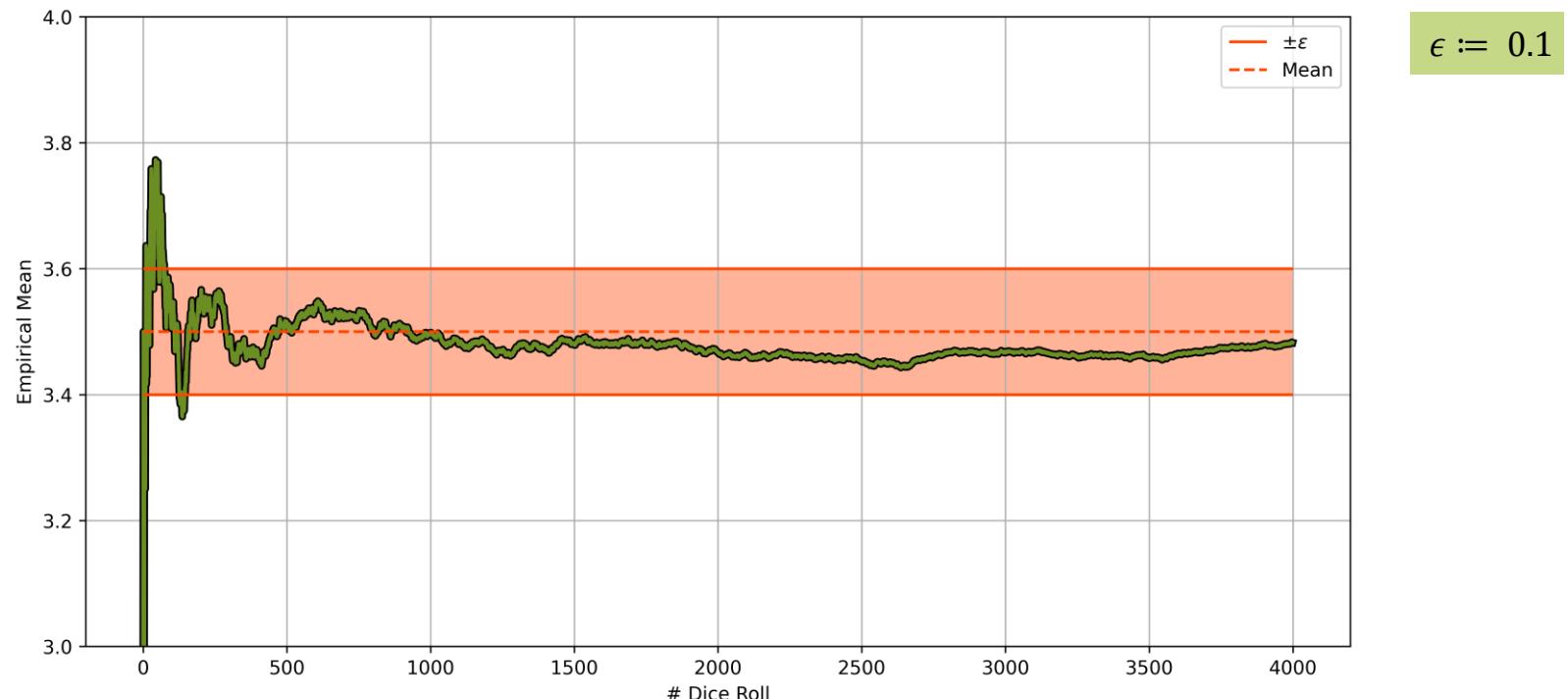
Why Tail Bounds

Empirical Mean of a Dice



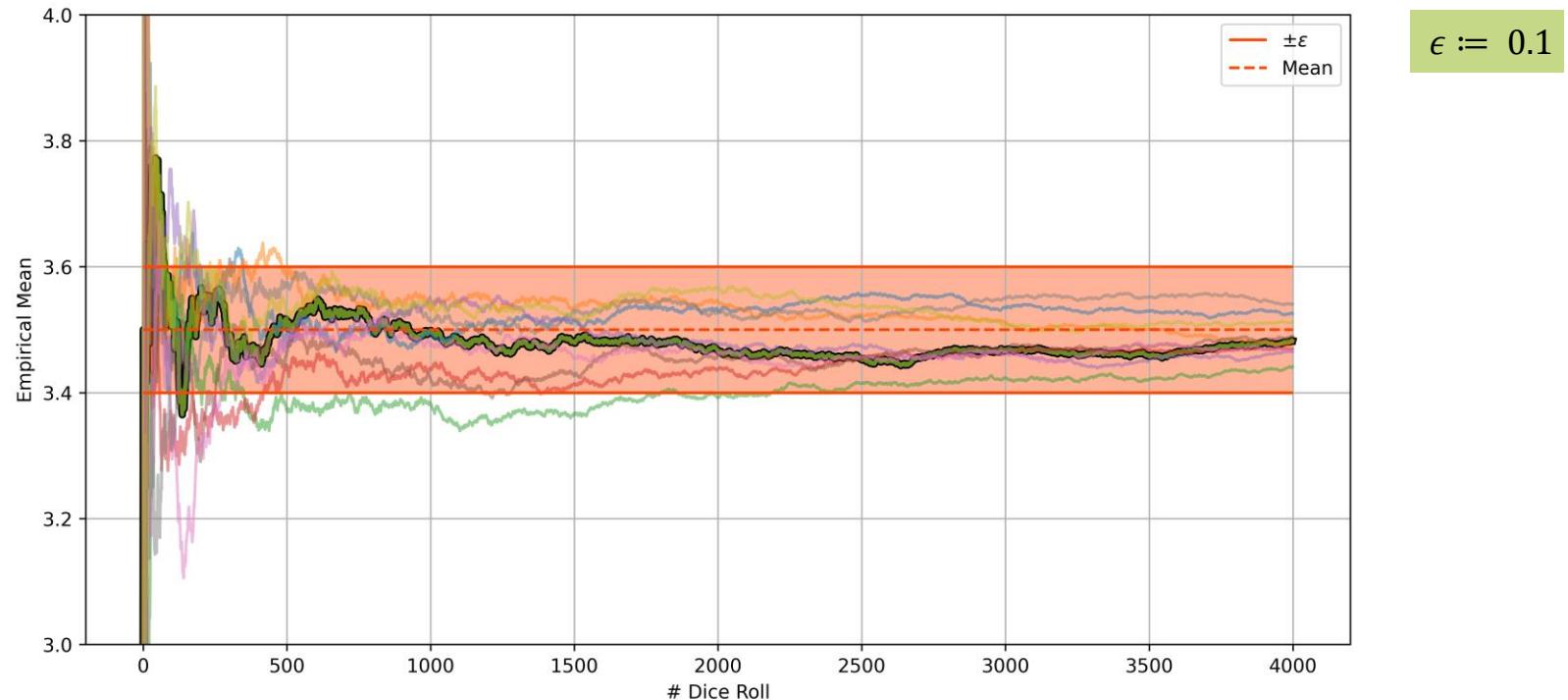
Why Tail Bounds

Empirical Mean of a Dice



Why Tail Bounds

Empirical Mean of a Dice



Höffding's Inequality

$$\mathbb{P}[|\bar{X}_t - \mu| \geq \epsilon] \leq 2 \exp\left[-\frac{2t\epsilon^2}{R^2}\right] := \delta$$

- Bounded form

$$|\bar{X}_t - \mu| \leq R \sqrt{\frac{\ln(2/\delta)}{2t}} := \epsilon$$

- Solve for t gives minimally required number of samples t_{min}

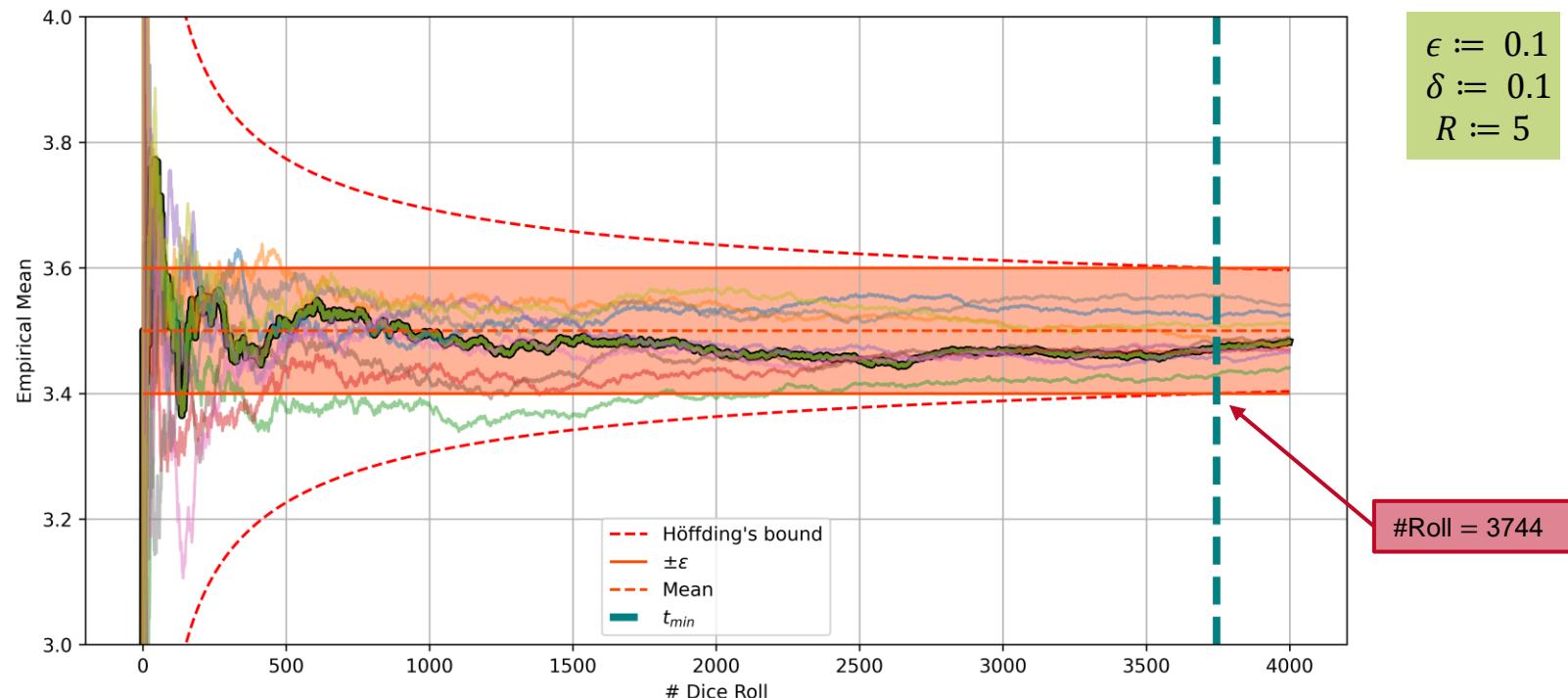
$$t_{min} = 2 \left(\frac{R}{\epsilon}\right)^2 \ln(2/\delta)$$

Independent and identically distributed random variable

- $X_1 \dots X_t \triangleq$ i.i.d random variables
- $a \leq X_i \leq b$
- $\epsilon \in \mathbb{R}^+$
- $R \triangleq$ Range

- $\bar{X}_t = \frac{1}{t} \sum_{i=1}^t X_i$

Empirical Mean of a Dice: Höffing's Bound



Tail Bounds

Bernstein's bound

$$\blacksquare |\bar{X}_t - \mu| \leq \sqrt{\frac{2 \Sigma^2 \ln(2/\delta)}{t}} + \frac{R \ln(2/\delta)}{3t}$$

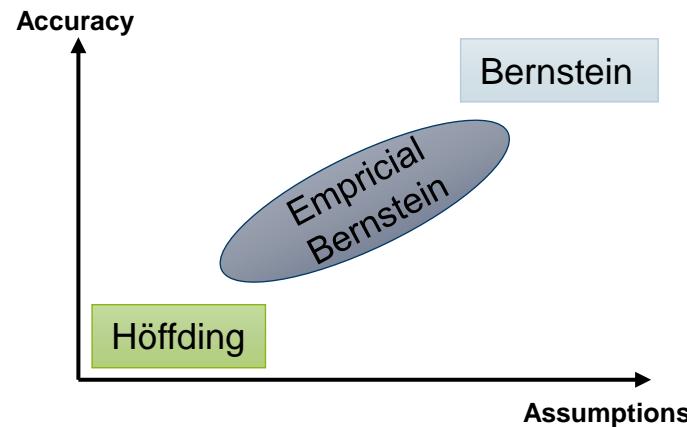
$\underbrace{\phantom{\sqrt{\frac{2 \Sigma^2 \ln(2/\delta)}{t}}}}_{\Theta(1/\sqrt{t})} \quad \underbrace{\phantom{\frac{R \ln(2/\delta)}{3t}}}_{\Theta(1/t)}$

For $\Sigma < R \rightarrow$ Bernstein “better”
For $\Sigma > R \rightarrow$ Höffding “better”

Höffding's bound

$$\blacksquare |\bar{X}_t - \mu| \leq R \sqrt{\frac{\ln(2/\delta)}{2t}}$$

$\underbrace{\phantom{\sqrt{\frac{\ln(2/\delta)}{2t}}}}_{\Theta(1/\sqrt{t})}$

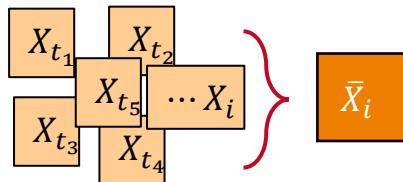


Empirical Bernstein stopping Improvements

Empirical Bernstein Stopping

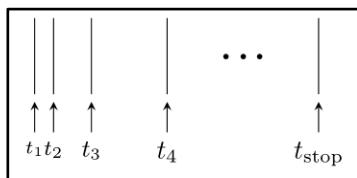
Pseudo Code Implementation: Improvements

- Batch the samples into one sample



- $X_1 \dots X_t \triangleq$ i.i.d random variables
- $\epsilon \in R^+$
- $1 - \delta \triangleq$ Confidence

- Update $c(t)$ with a growing gap



- $R \triangleq$ Range
- $\bar{X}_t = \frac{1}{z} \sum_{i=1}^t X_i$
- $c(t) = \sqrt{\frac{2\bar{V}_t \ln(3/\delta_t)}{t}} + \frac{3R \ln(3/\delta_t)}{t}$